Scheduling in Windows

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These materials are part of the *Windows Operating System Internals Curriculum Development Kit*, developed by David A. Solomon and Mark E. Russinovich with Andreas Polze

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“Although programs and processes appear similar on the surface, they are fundamentally different. A program is a static sequence of instructions, whereas a process is a container for a set of resources used when executing the instance of the program. At the highest level of abstraction, a Windows process comprises the following:

- A private virtual address space, which is a set of virtual memory addresses that the process can use
- An executable program, which defines initial code and data and is mapped into the process’s virtual address space
- A list of open handles to various system resources, such as semaphores, communication ports, and files, that are accessible to all threads in the process
- A security context called an access token that identifies the user, security groups, and privileges associated with the process
- A unique identifier called a process ID (internally called a client ID)
- At least one thread of execution”

“A thread is the entity within a process that Windows schedules for execution. Without it, the process’s program can’t run. Although threads have their own execution context, every thread within a process shares the process’s virtual address space (in addition to the rest of the resources belonging to the process), meaning that all the threads in a process can write to and read from each other’s memory. Threads cannot accidentally reference the address space of another process, however, unless the other process makes available part of its private address space as a shared memory section (called a file mapping object in the Windows API) or unless one process has the right to open another process to use cross-process memory functions such as ReadProcessMemory and WriteProcessMemory.”
Principles of Windows scheduling

- Preemptive scheduler (both kernel and user!)
- 32 priority levels
  - (One of the) Thread with the highest priority runs
  - Round robin between threads with same priority
- Threads run for a fixed time (quantum)
- No central scheduler, scheduling is driven by events
- Priority of threads can change runtime
Ready: A thread in the ready state is waiting to execute. When looking for a thread to execute, the dispatcher considers only the pool of threads in the ready state.

- Standby: A thread in the standby state has been selected to run next on a particular processor. When the correct conditions exist, the dispatcher performs a context switch to this thread. Only one thread can be in the standby state for each processor on the system. Note that a thread can be preempted out of the standby state before it ever executes (if, for example, a higher priority thread becomes runnable before the standby thread begins execution).

- Running: Once the dispatcher performs a context switch to a thread, the thread enters the running state and executes. The thread’s execution continues until its quantum ends (and another thread at the same priority is ready to run), it is preempted by a higher priority thread, it terminates, it yields execution, or it voluntarily enters the wait state.

- Waiting: A thread can enter the wait state in several ways: a thread can voluntarily wait for an object to synchronize its execution, the operating system can wait on the thread’s behalf (such as to resolve a paging I/O), or an environment subsystem can direct the thread to suspend itself. When the thread’s wait ends, depending on the priority, the thread either begins running immediately or is moved back to the ready state.

- Transition: A thread enters the transition state if it is ready for execution but its kernel stack is paged out of memory. Once its kernel stack is brought back into memory, the thread enters the ready state.

- Terminated: When a thread finishes executing, it enters the terminated state. Once the thread is terminated, the executive thread block (the data structure in nonpaged pool that describes the thread) might or might not be deallocated. (The object manager sets policy regarding when to delete the object.)

- Initialized: This state is used internally while a thread is being created.
Priority levels (kernel)

- **31**: NOT hard/soft real-time, just priority is constant
- **16 “real-time”**: Priority of the thread can be changed
- **15 dynamic**: If nothing else to run, counts spare cycles
## Priority levels (Windows API, GUI)

<table>
<thead>
<tr>
<th>Name of priority levels</th>
<th>Value of priority</th>
</tr>
</thead>
<tbody>
<tr>
<td>Realtime</td>
<td>31</td>
</tr>
<tr>
<td>High</td>
<td>16</td>
</tr>
<tr>
<td>Above Normal</td>
<td>15</td>
</tr>
<tr>
<td>Normal</td>
<td></td>
</tr>
<tr>
<td>Below Normal</td>
<td></td>
</tr>
<tr>
<td>Idle</td>
<td>1</td>
</tr>
</tbody>
</table>
Windows API vs. kernel names

- Threads: 7 different relative priority
- Mapping:

<table>
<thead>
<tr>
<th>Win32 thread priorities</th>
<th>Win32 process priority levels</th>
</tr>
</thead>
<tbody>
<tr>
<td></td>
<td>Realtime</td>
</tr>
<tr>
<td>Time-critical</td>
<td>31</td>
</tr>
<tr>
<td>Highest</td>
<td>26</td>
</tr>
<tr>
<td>Above-normal</td>
<td>25</td>
</tr>
<tr>
<td>Normal</td>
<td>24</td>
</tr>
<tr>
<td>Below-normal</td>
<td>23</td>
</tr>
<tr>
<td>Lowest</td>
<td>22</td>
</tr>
<tr>
<td>Idle</td>
<td>16</td>
</tr>
</tbody>
</table>
I/O priority

- Since Vista
- 5 different I/O priority for requests, e.g.
  - Critical: Dirty page writer
  - Low: Desktop search indexer
- I/O bandwidth allocation
DEMO  Changing priority

- Task manager

![Task Manager Screen Shot](image)

![Properties Dialog](image)
From Windows Internals curriculum:
“The dispatcher ready queues (KiDispatcherReadyListHead) contain the threads that are in the ready state, waiting to be scheduled for execution. There is one queue for each of the 32 priority levels. To speed up the selection of which thread to run or preempt, Windows maintains a 32-bit bit mask called the ready summary (KiReadySummary). Each bit set indicates one or more threads in the ready queue for that priority level. (Bit 0 represents priority 0, and so on.)”

The length of the clock interval varies according to the hardware platform. The frequency of the clock interrupts is up to the HAL, not the kernel. For example, the clock interval for most x86 uniprocessors is about 10 milliseconds and for most x86 and x64 multiprocessors it is about 15 milliseconds.
On Windows Vista, threads run by default for 2 clock intervals; on Windows Server systems, by default, a thread runs for 12 clock intervals. The rationale for the longer default value on server systems is to minimize context switching. By having a longer quantum, server applications that wake up as the result of a client request have a better chance of completing the request and going back into a wait state before their quantum ends.

Threads in the foreground process run with a quantum of 6 clock ticks, whereas threads in other processes have the default workstation quantum of 2 clock ticks. In this way, when you switch away from a CPU-intensive process, the new foreground process will get proportionally more of the CPU, because when its threads run they will have a longer turn that background threads (again, assuming the thread priorities are the same in both the foreground and background processes).

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**From Windows Internals Curriculum**

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**Short or Long, Variable or Fixed:**

<table>
<thead>
<tr>
<th>HKLM\SYSTEM\CurrentControlSet\Control\PriorityControl\Win32PrioritySeparation</th>
</tr>
</thead>
</table>
Length of the quantum

- Clockres.exe utility
  - Length of Clock tick
- Adjust quantum length
- Perfmon
- Windows Performance Analyzer
  - Timeline by Process, Thread view
Contents

- Windows scheduling (basics)
- Windows 8: Windows Store applications
- Windows scheduling (advanced)
Windows 8: new application model

- Design goals:
  - Low powers and resource-consumption
  - Easier deployment and upgrade
  - Separated applications (security, reliability)
  - ...

- Solution:
  - New API: WinRT
  - App store: Windows Store
  - New application lifecycle
  - ...
Windows Store application lifecycle

If saved state exists, can be restore; otherwise start from scratch

Use changes to other app. OS suspends the process (5 sec). App can change its state.

Activated ➔ Running ➔ Suspending

NotRunning ➔ Running ➔ Suspending

Resuming ➔ Suspending

Terminating

No CPU time OS can terminate app

DEMO  Windows Store applications

- Windows Store
- Changing between Windows Store apps
- Process Explorer: Suspended state
Contents

- Windows scheduling (basics)
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• Five types:
  I/O completion
  Wait completion on events or semaphores
  When threads in the foreground process complete a wait
  When GUI threads wake up for windows input
  For CPU starvation avoidance

• Quantum decremented by 1 when you come out of a wait
  So that threads that get boosted after I/O completion won't keep running and never experiencing quantum end
  Prevents I/O bound threads from getting unfair preference over CPU bound threads
Preventing starvation

- OS checks runnable threads (every 1 sec)

- If a thread has not been executed since 300 sec
  - change priority to 15 (max in dynamic),
  - increase quantum,
  - for 1 quantum.
Symmetric Multiprocessing (SMP)

- All CPU is equal
  - Shared address space
  - Interrupts can be served by an CPU

- Implementation limit (length of a bit vector):
  - 32 CPUs on 32 bit systems
  - 64 CPUs on 64 bit systems

- Change in Windows 7 / Server 2008 R2
  - Groups of logical processors
  - Supporting 4 * 64 CPU
  - NUMA support
Multiprocessor scheduling

- Threads can run on any CPU by default, but
  - OS tries to keep on CPU, where it run (“soft affinity”)
  - Can be set to use only selected CPUs (“hard affinity”)

- No “master processor”

- Dispatcher queues:
  - Before Windows Server 2003: one global queue
  - Windows Server 2003: per CPU queues
Affinity is a bit mask where each bit corresponds to a CPU number

• Hard Affinity specifies where a thread is permitted to run
  • Defaults to all CPUs
  • Thread affinity mask must be subset of process affinity mask, which in turn must be a subset of the active processor mask

Functions to change: SetThreadAffinityMask, SetProcessAffinityMask, SetInformationJobObject
Windows 7 changes

- Core Parking (server)
  - Use fewer processor cores
  - Not used cores going to standby
- Time coalescing
  - Timers with same periodicity are merged
- Dynamic Fair Share Scheduling (DFSS)
  - for Remote Desktop
  - Every session gets a share
  - If share is exhausted, thread cannot run
- Eliminating global locks in scheduler
Summary

- Process ↔ Thread

- Scheduling:
  - Priority levels
  - Round robin / quantum